

Homework Assignment #1 (100 points, weight 5%)

Due: Thursday Feb 9, at 1:00 p.m. (in lecture);

assignments with lateness between 1min-24hs will have a discount of 10%; after 24hs, not accepted; please drop off late assignments under my office door (STE5027).

Propositional Logic

1. (12 points) Use logical equivalences to show that $[(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)] \rightarrow r$ is a tautology.
2. (12 points; each 2+2+2 points=truth table+DNF+CNF)
For each of the following compound propositions give its truth table and derive an equivalent compound proposition in disjunctive normal form (DNF) and in conjunctive normal form (CNF).
 - (a) $(p \rightarrow q) \rightarrow r$
 - (b) $(p \wedge \neg q) \vee (p \leftrightarrow r)$

Predicate Logic

3. (15 points) For each of the given statements:
 - 1 - Express each of the statements using quantifiers and propositional functions.
 - 2 - Form the negation of the statement so that no negation is to the left of the quantifier.
 - 3 - Express the negation in simple English. (Do not simply use the words “it is not the case that...”).
 - (a) Some drivers do not obey the speed limit.
 - (b) All Swedish movies are serious.
 - (c) No one can keep a secret.
 - (d) No monkey can speak French.
 - (e) There is someone in the class who does not have a good attitude.
4. (10 points) Translate these system specifications into English where the predicate $S(x, y)$ is “ x is in state y ” and where the domain for x and y consists of all systems and all possible states, respectively.
 - (a) $\exists S(x, \text{open})$
 - (b) $\forall x(S(x, \text{malfunctioning}) \vee S(x, \text{diagnostic}))$
 - (c) $\exists x S(x, \text{open}) \vee \exists x S(x, \text{diagnostic})$

- (d) $\exists x \neg S(x, \text{available})$
 (e) $\forall x \neg S(x, \text{working})$
5. (3+3+3+3=12 marks) Rewrite the following statements so that all negation symbols immediately precede predicates (that is, no negation is outside a quantifier or an expression involving logical connectives). Show all the steps in your derivation.
- (a) $\neg \forall x \exists y P(x, y)$
 (b) $\neg \exists y (Q(y) \wedge \forall x \neg R(x, y))$
 (c) $\neg \exists y (\exists x R(x, y) \vee \forall x S(x, y))$
 (d) $\neg \exists y (\forall x \exists z T(x, y, z) \vee \exists x \forall z U(x, y, z))$
6. (10 points) Prove these logical equivalences, assuming that the domain is nonempty. You will probably have to use a proof by cases on the two possible values of proposition $\forall y Q(y)$ and $\exists y Q(y)$ respectively. This proof will use word arguments (not symbolic formula manipulation).
- (a) $\forall x (\forall y Q(y) \rightarrow P(x)) \equiv \forall y Q(y) \rightarrow \forall x P(x)$
 (b) $\exists x (\exists y Q(y) \rightarrow P(x)) \equiv \exists y Q(y) \rightarrow \exists x P(x)$
7. (10 points) A statement is in prenex normal form (PNF) if and only if all quantifiers occur at the beginning of the statement (without negations), followed by a predicate involving no quantifiers. Put the following statement in prenex normal form:
 (Hint: your first step should rename one of the two x 's as y ; check useful valid equivalences in exercises 48,49 page 62 of 6th edition)
- (a) $\exists x P(x) \vee \exists x Q(x) \vee A$, where A is a proposition not involving any quantifiers.
 (b) $\exists x P(x) \rightarrow \exists x Q(x)$

Rules of Inference

8. (9 points) For each of these arguments, determine whether the argument is correct or incorrect and explain why.
- (a) Everyone born in Ottawa has eaten a beaver tail. Susan has never eaten a beaver tail. Therefore Susan was not born in Ottawa.
 (b) A convertible car is fun to drive. Joe's car is not a convertible. Therefore, Joe's car is not fun to drive.
 (c) Emma likes all fine restaurants. Emma likes the restaurant "Le Cordon Bleu". Therefore, "Le Cordon Bleu" is a fine restaurant.

9. (10 points) Give a formal proof, using known rules of inference, to establish the conclusion of the argument (3rd statement) using the first 2 statements as premises, where the domain of all quantifiers is the same.

Remember that a formal proof is a sequence of steps, each with a reason noted beside it; each step is either a premise, or is obtained from previous steps using inference rules.

- premise: $\forall x(P(x) \vee Q(x))$
- premise: $\forall x((\neg P(x) \wedge Q(x)) \rightarrow R(x))$
- conclusion: $\forall x(\neg R(x) \rightarrow P(x))$